

On the Jones polynomials of checkerboard colorable virtual knots

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Abstract

In this paper we study the Jones polynomials of virtual links and abstract links. It is proved that a certain property of the Jones polynomials of classical links is valid for virtual links which admit checkerboard colorings.

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1 Introduction

In 1996, L. H. Kauffman introduced the notion of a virtual knot, which is motivated by study of knots in a thickened surface and abstract Gauss codes, cf. [8, 9]. M. Goussarov, M. Polyak, and O. Viro [1] proved that the natural map from the category of classical knots to the category of virtual knots is injective; namely, if two classical knot diagrams are equivalent as virtual knots, then they are equivalent as classical knots. Thus, virtual knot theory is a generalization of knot theory. It is also found in their paper [1] that the notion of a virtual knot is helpful to study of finite type invariants.

Kauffman defined the Jones polynomial of a virtual knot, which is also called the normalized bracket polynomial or the f -polynomial (cf. [9]). In this paper, according to [9], we call it the f -polynomial instead of the Jones polynomial, since the definition is different from Jones' in [2, 3]. Finite type invariants derived from the f -polynomials are studied in [9], and it is proved that a certain property of them (Corollary 14 of [9]) is hold in the category of virtual knots.

The f -polynomial (Jones polynomial) of a virtual link is quite different from f -polynomials of classical links. For a Laurent polynomial f on valuable A , we denote by $\text{EXP}(f)$ the set of integers appearing as exponents of f . For example, if $f = 3A^{-2} + 6A - 7A^5$, then $\text{EXP}(f) = \{-2, 1, 5\}$. It is well-known

that for a classical link L with n components, the f -polynomial satisfies that $\text{EXP}(f) \subset 4\mathbf{Z}$ if n is odd, and $\text{EXP}(f) \subset 4\mathbf{Z} + 2$ if n is even. However, this is not true for a virtual knot/link in general. In this paper we introduce the notion of *checkerboard coloring* of a virtual link diagram as a generalization of checkerboard coloring of a classical link diagram.

Theorem 1 *Let f be the f -polynomial of a virtual link L with n components. Suppose that L has a virtual link diagram which admits a checkerboard coloring. Then $\text{EXP}(f) \subset 4\mathbf{Z}$ if n is odd, and $\text{EXP}(f) \subset 4\mathbf{Z} + 2$ if n is even.*

For example the virtual knot diagram illustrated in Figure 1 (a) admits a checkerboard coloring and the f -polynomial is $A^4 + A^{12} - A^{16}$. So $\text{EXP}(f) \subset 4\mathbf{Z}$. On the other hand, virtual knot diagram illustrated in Figure 1 (b) does not admit a checkerboard coloring and the f -polynomial is $-A^{10} + A^6 + A^4$. Theorem 1 implies that this diagram is never equivalent to a diagram that admits a checkerboard coloring.

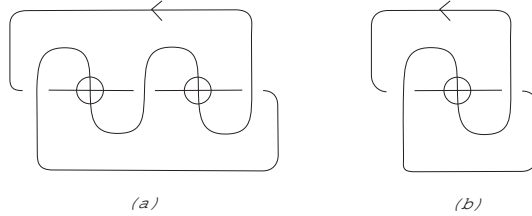


Figure 1:

If a virtual link diagram is alternating (the definition is given later), then the diagram admits a checkerboard coloring. Thus we have the following.

Corollary 2 *Let f be the f -polynomial of a virtual link L with n components. Suppose that L has an alternating virtual link diagram. Then $\text{EXP}(f) \subset 4\mathbf{Z}$ if n is odd, and $\text{EXP}(f) \subset 4\mathbf{Z} + 2$ if n is even.*

By this corollary, we see that the virtual knot represented by Figure 1 (b) is not equivalent to an alternating diagram.

2 Virtual link diagram and abstract link diagram

A *virtual link diagram* is a closed oriented 1-manifold generically immersed in \mathbf{R}^2 such that each double point has information of a crossing (as in classical

knot theory) or a virtual crossing which is indicated by a small circle around the double point. The moves of virtual link diagrams illustrated in Figure 2 are called *generalized Reidemeister moves*. Two virtual link diagrams are said to be *equivalent* if they are related by a finite sequence of generalized Reidemeister moves. We call the equivalence class of a virtual link diagram a *virtual link*.

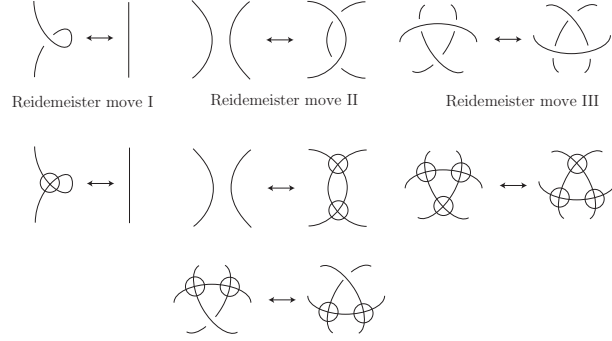


Figure 2:

A pair $P = (\Sigma, D)$ of a compact oriented surface Σ and a link diagram D in Σ is called an *abstract link diagram* (ALD) if $|D|$ is a deformation retract of Σ , where $|D|$ is a graph obtained from D by replacing each real/virtual crossing point with a vertex. For an ALD, $P = (\Sigma, D)$, if there is an orientation preserving embedding $f : \Sigma \rightarrow F$ into a closed oriented surface F , $f(D)$ is a link diagram in F . We call it a *link diagram realization* of P in F . In Figure 3, we show two abstract link diagrams and their link diagram realizations. Two ALDs $P = (\Sigma, D)$, $P' = (\Sigma', D')$ are related by an *abstract Reidemeister move* (of type I, II or III) if there is a closed oriented surface F and link diagram realizations of P and P' in F which are related by a Reidemeister move (of type I, II or III) in F . Two ALDs are *equivalent* if they are related by a finite sequence of abstract Reidemeister moves. We call the equivalence class of an ALD an *abstract link*.

In [6] a map

$$\phi : \{\text{virtual link diagrams}\} \longrightarrow \{\text{ALDs}\}$$

was defined. The idea of this map is illustrated in Figure 4. Refer to [6] for the definition. We call $\phi(D)$ an *ALD associated with a virtual link diagram* D . The ALDs in Figure 3 (a) and (b) are ALDs associated with the virtual link diagrams in Figure 1 (a) and (b) respectively.

Theorem 3 ([6]) *The map ϕ induces a bijection*

$$\Phi : \{\text{virtual links}\} \longrightarrow \{\text{abstract links}\}$$

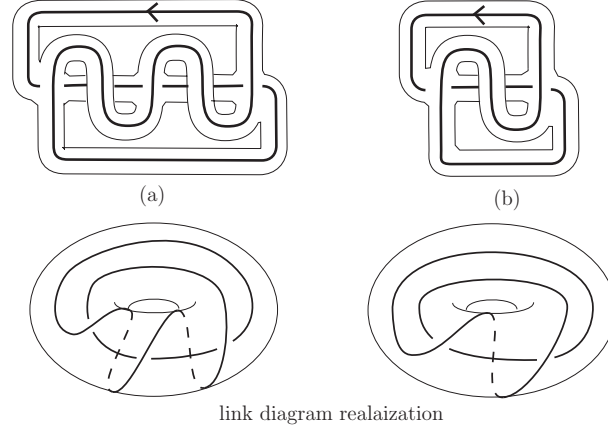


Figure 3:



Figure 4:

Let $P = (\Sigma, D)$ be a pair of a compact oriented surface Σ and a link diagram D in Σ . A *checkerboard coloring* is a coloring of the all components of $\Sigma - |D|$ by two colors, say black and white, such that two components of $\Sigma - |D|$ which are adjacent by an edge of D have always distinct colors.

We say that a virtual link diagram *admits a checkerboard coloring* or it is *checkerboard colorable* if the associated ALD admits a checkerboard coloring.

3 The f -polynomials of abstract link diagrams

An ALD, $P = (\Sigma, D)$, is said to be *unoriented* if the diagram D is unoriented. There is a unique map

$$\langle \rangle: \{\text{unoriented ALDs}\} \longrightarrow \Lambda = \mathbf{Z}[A, A^{-1}]$$

satisfying the following rules.

- (i) $\langle T \rangle_F = 1$ where T is a one-component trivial ALD,

(ii) $\langle T \amalg D \rangle = (-A^2 - A^{-2}) \langle D \rangle$ if D is not empty, where \amalg means the disjoint union, and

$$(iii) \langle \begin{array}{c} \diagup \diagdown \\ \diagdown \diagup \end{array} \rangle = A \langle \begin{array}{c} \frown \\ \smile \end{array} \rangle + A^{-1} \langle \begin{array}{c} \smile \\ \frown \end{array} \rangle.$$

Then $\langle \rangle$ is an invariant under abstract Reidemeister moves II and III. We call it the *Kauffman bracket polynomial* of ALD, cf. [4].

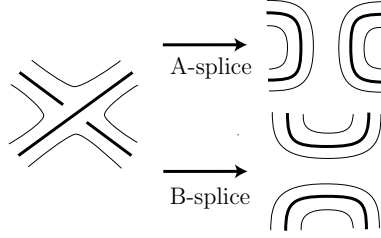


Figure 5:

Let $P = (\Sigma, D)$ be an unoriented ALD. Replacing the neighborhood of a double point as in Figure 5, we have another unoriented ALD. We call it an unoriented ALD obtained from D by doing an *A-splice* or *B-splice* at the crossing point. An unoriented trivial ALD obtained from P by doing an A-splice or B-splice at each crossing point is said to be a *state* of P . From the definition of $\langle \rangle$, we see

$$\langle P \rangle = \sum_S A^{\natural(S)} (-A^2 - A^{-2})^{\sharp(S)-1},$$

where S runs over all of states of D , $\natural(S)$ is the number of A-splice minus that of B-splice used for obtaining S and $\sharp(S)$ is the number of components of S .

For an ALD, $P = (\Sigma, D)$, the writhe $\omega(P)$ is defined by the number of positive crossings minus the number of negative crossings. Then we define the *normalized bracket polynomial* or the *f-polynomial* of P by

$$f_P(A) = (-A^3)^{-\omega(P)} \langle P \rangle.$$

By normalizing by $(-A^3)^{-\omega(P)}$, this value is preserved under abstract Reidemeister moves of type I. Thus this is an invariant of an abstract link. This invariant was defined in [4], where it is called the Jones polynomial of P . It should be noted that the bijection Φ preserves the *f-polynomial*.

4 Proof of Theorem 1

Let p be a crossing point of an ALD, $P = (\Sigma, D)$. Let $P_0 = (\Sigma_0, D_0)$ and $P_\infty = (\Sigma_\infty, D_\infty)$ be ALDs obtained from P by splicing at p orientation coherently and orientation incoherently, respectively. Note that D_∞ does not inherit an orientation from D . The crossing point p is either (i) a self-intersection of an immersed loop of D or (ii) an intersection of two immersed loops. Let α and α' be the immersed open arcs obtained from the loop (in case (i)) or from the two loops (in case (ii)) by cutting at p . Choose one of them, say α , and we give an orientation to D_∞ which is induced from that of D except α (and hence the orientation is reversed on α). Let C be the set of crossing points of D , except p , such that the sign of the crossing point does not change in D and D_∞ ; in other word, at each crossing point belonging to C , both of the two intersecting arcs are contained in $D - \alpha$ or both of them are in α . Let C' be the set of crossing points of D , except p , such that the sign of the crossing point changes in D and D_∞ ; in other word, at each crossing point belonging to C' , one of the two intersecting arcs is contained in $D - \alpha$ and the other is in α . Let k (or ℓ , resp.) be the number of positive crossings of C (resp. C') minus the number of negative crossings of C (resp. C').

Lemma 4 *In the above situation, let f , f_0 and f_∞ be the f -polynomials of P , P_0 and P_∞ , respectively. Then we have*

$$f = \begin{cases} -A^{-2}f_0 - (-A^3)^{-2\ell}A^{-4}f_\infty, & \text{if } p \text{ is a positive crossing,} \\ -A^{+2}f_0 - (-A^3)^{-2\ell}A^{+4}f_\infty, & \text{if } p \text{ is a negative crossing.} \end{cases}$$

Proof. If p is a positive crossing, then the writhes are $\omega(D) = k + \ell + 1$, $\omega(D_0) = k + \ell$ and $\omega(D_\infty) = k - \ell$. Since $\langle P \rangle = A \langle P_0 \rangle + A^{-1} \langle P_\infty \rangle$, we have the result. The case that p is a negative crossing is similar. \square

Remark. In the remark of Section 5 of [9](page 677), an equation which is similar to Lemma 4 is given. However, it seems to be forgotten there to take account of the term $(-A^3)^{-2\ell}$. In consequence, the recursion formula of Theorem 13 of [9] is as follows:

$$v_n(G_*) = \sum_{k=0}^{n-1} \frac{2^{n-k}}{(n-k)!} \{ (1 - (-1)^{n-k})v_k(G_0) + \{ (2-3\ell)^{n-k} - (-2-3\ell)^{n-k} \} v_k(G_\infty) \}.$$

By this formula, Corollary 14 of [9] is still true.

Corollary 5 (cf. Theorem 13 of [9]) *Let f be the f -polynomial of an ALD with n components. Then $f(1) = (-2)^{n-1}$. In particular, f -polynomials of ALDs are not zero.*

Proof. It follows from Lemma 4 by induction on the number of (real) crossing points. \square

Since Φ preserves the f -polynomials, Theorem 1 is equivalent to the following theorem.

Theorem 6 *Let f be the f -polynomial of an ALD, $P = (\Sigma, D)$, with n components. Suppose that P admits a checkerboard coloring. Then $\text{EXP}(f) \subset 4\mathbf{Z}$ if n is odd, and $\text{EXP}(f) \subset 4\mathbf{Z} + 2$ if n is even.*

Proof. For a state S of P , we define $I(S)$ by

$$I(S) = A^{\sharp(S)}(-A^2 - A^{-2})^{\sharp(S)-1}$$

so that the bracket polynomial of P is the sum of $I(S)$ for all states S . Let $\text{ind}(S)$ be a value in $\mathbf{Z}_4 = \{0, 1, 2, 3\}$ such that $I(S) \subset 4\mathbf{Z} + \text{ind}(S)$.

Every state of P has a unique checkerboard coloring induced from the checkerboard coloring of P , see Figure 6. (Figure 7 is an example of an ALD with a checkerboard coloring and a state with the induced checkerboard coloring.) Using this fact, we prove that $\text{ind}(S) = \text{ind}(S')$ for any states S and S' of P . It is sufficient to prove this in a special case that S and S' are the same state except a crossing point, say p , of D where S and S' are as in Figure 8. For this state S , there are two cases (A) and (B) as in Figure 9. The case (C) does not occur, because a state as in (C) does not have a checkerboard coloring induced from the checkerboard coloring of P . In both cases (A) and (B), we have $I(S') = A^{\sharp(S) \pm 2}(-A^2 - A^{-2})^{\sharp(S)-1 \pm 1}$ and $\text{ind}(S) = \text{ind}(S')$.

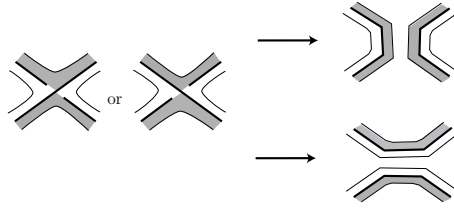


Figure 6:

Now we have that $\text{EXP}(f) \subset 4\mathbf{Z} + i$ where $i = \text{ind}(S)$ for any state S of P . We denote this number i by $\text{ind}(f)$. The remaining task is to prove this index is 0 if n is odd, and 2 if n is even. This is proved by induction on the number of (real) crossing points of P . If P has no real crossing points, then this is obvious by the definition of the f -polynomial. If there is a crossing point, say p , apply Lemma 4. Note that P_0 and P_∞ have checkerboard colorings, and $\text{EXP}(f_0) \subset 4\mathbf{Z} + \text{ind}(f_0)$ and $\text{EXP}(f_\infty) \subset 4\mathbf{Z} + \text{ind}(f_\infty)$. Since $f \neq 0$ and $f_0 \neq 0$ (Corollary 5), it follows

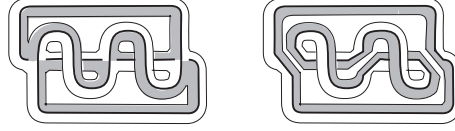


Figure 7:

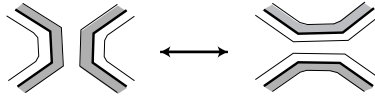


Figure 8:

from the equation in Lemma 4 that $\text{ind}(f) = \text{ind}(f_0) + 2 \in \mathbf{Z}_4$. The ALD P_0 has fewer crossing points than P and has a checkerboard coloring. By induction hypothesis, $\text{ind}(f_0)$ is 0 if n' is odd, and 2 if n' is even, where n' is the number of components of P_0 . Since $n' = n \pm 1$, we have that $\text{ind}(f)$ is 0 if n is odd, and 2 if n is even. \square

5 Alternating virtual link diagrams and ALDs

An ALD or a virtual link diagram is *alternating* if we meet over and under crossing points alternatively when we travel along each component of the diagram twice.

Lemma 7 *For an ALD, $P = (\Sigma, D)$, the following conditions are equivalent.*

- (i) *By applying crossing changes, P changes into an alternating ALD.*
- (ii) *P has a checkerboard coloring.*

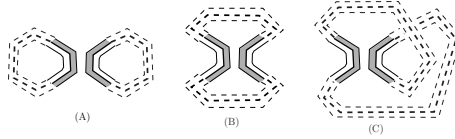


Figure 9:

Proof of Lemma 7. If P has a checkerboard coloring, change each real crossing according to the coloring as in the most left figure of Figure 6. Conversely if P is an alternating ALD, then give a checkerboard coloring near each crossing point as in the picture, which is extended to a checkerboard coloring of P . \square

Proof of Corollary 2. It follows from Theorem 1 and Lemma 7. \square

Remark. M. B. Thistlethwaite [11] and K. Murasugi [10] showed that the f -polynomial (Jones polynomial) of a non-split alternating link is alternating, namely it is in a form of $A^\alpha \sum c_i A^{4i}$ such that $c_i c_j \geq 0$ for $i \equiv j \pmod{2}$ and $c_i c_j \leq 0$ for $i \not\equiv j \pmod{2}$. This does not hold in virtual knot theory. The f -polynomial of a virtual knot in Figure 10 is $A^{12} + 3A^{16} - 4A^{20} + 3A^{24} - 4A^{28} + 4A^{32} - 3A^{36} + A^{40}$.

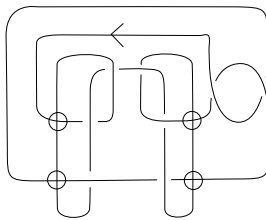


Figure 10:

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